

PATHOS VERTEX SEMI-MIDDLE GRAPH OF A TREE

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Abstract: In this communication, the pathos vertex semi-middle graph of a tree is introduced. Its study is concentrated only on trees. We present a characterization of those graphs whose pathos vertex semi-middle graph of a tree is planar, crossing number one and crossing number two. Also we establish a characterization of graphs whose pathos vertex semi-middle graph of a tree is noneulerian, hamiltonian.

Keywords and Phrases: Crossing number; Middle graph; Planar; Semientire graph.

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1. Introduction

All graphs considered here are finite, undirected without loops or multiple edges. We refer the terminology of [1].

The concept of pathos of a graph G was introduced by Harary [2], as a collection of minimum number of edge disjoint open paths whose union is G . The path number of a graph G is the number of paths in pathos.

A graph is said to be embedded in a surface S when it is drawn on S so that no two edges intersect. A graph is planar if it can be embedded in the plane.

The crossing number $C_r(G)$ of G is the least number of intersection of pairs of edges in any embedding of G in the plane. Obviously G is planar if and only if $C_r(G) = 0$.

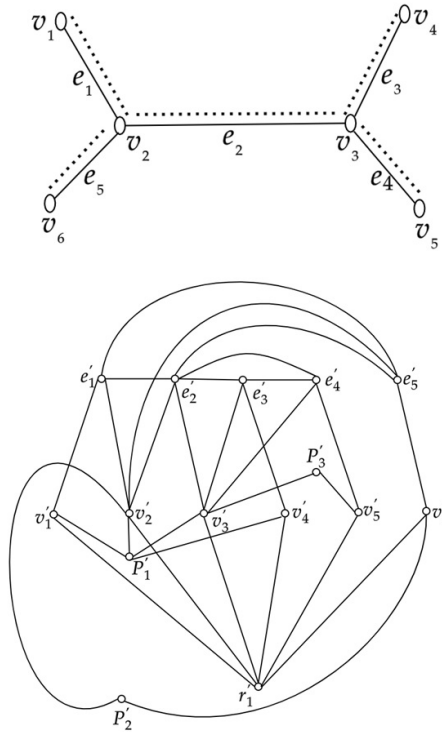


Figure 1. T and $PM_v(T)$

The *edgedegree* [5] of an edge uv of a tree T is the sum of the degrees of u and v . The pathlength is the number of edges that lie on a particular path P_i of pathos of T . A pendent pathos is a path P_i of pathos having unit length, which corresponds to a pendent edge in T .

In the paper [4] and [6] defined the concept of on pathos litact graph of a tree and pathos cutvertex jump graph of a tree. In this context, the present work focus to define the pathos vertex semi-middle graph of a tree. The pathos vertex semi-middle graph of a tree denoted by $PM_v(T)$ is the graph whose vertex set is $V(T)UE(T)UR(T)UP_i(T)$ and two vertices of $PM_v(T)$ are adjacent if and only if they corresponds to two adjacent edges of T or one corresponds to a vertex and other to an edge incident with it or one corresponds to a vertex other to a region in which vertex lies on the region or one corresponds to a vertex and other to a path of pathos in which vertex lies on the path of pathos since the system of pathos for a tree is not unique, the corresponding pathos vertex semi-middle graph of a tree is also not unique. The tree T and its pathos vertex semi-middle graph of a tree $PM_v(T)$ is depicted in the figure 1.

2. Preliminaries.

The following results will be useful in our results.

Theorem 2.1 [3]. *For any graph G , $M_v(G)$ is always non-separable.*

Theorem 2.2 [3]. *For any graph G , p vertices, q edges and l regions then $M_v(G)$ has $(p + q + r)$ vertices and $q + \sum_{i=1}^q \frac{1}{2}\{d(e_i)\} + \sum_{j=1}^r d(r_j)$ edges. Where $d(e_i)$ is the edgedegree of a edge e_i and $d(r_j)$ is the degree of a region r_j .*

Theorem 2.3 [3]. *For any graph G , $M_v(G)$ is planar if and only if G is a path.*

Theorem 2.4 [1]. *A finite graph G is Eulerian if and only if all its vertex degree are even.*

3. Pathos Vertex Semi-Middle Graph of a Tree.

We begin with the following observation.

Observation 3.1. *If v is a pendant vertex of a tree T , then the degree of a corresponding vertex v' in $PM_v(T)$ is odd.*

Theorem 3.1. *For any tree T , $PM_v(T)$ is always non-separable.*

Proof. We establish the following cases.

Case 1. Suppose T be any tree. Let $v'_1, v'_2, v'_3, \dots, v'_n$ be the vertices of $PM_v(T)$ corresponds to the vertices $v_1, v_2, v_3, \dots, v_n$ of T and $e'_1, e'_2, e'_3, \dots, e'_{n-1}$ be the vertices of $PM_v(T)$ corresponds to the edges $e_1, e_2, e_3, \dots, e_{n-1}$ of T . Let $P'_1, P'_2, P'_3, \dots, P'_n$ be the pathosvertices of $PM_v(T)$ corresponds to the path $P_1, P_2, P_3, \dots, P_n$. By Theorem 2.1, $M_v(G)$ is non-separable. Also in $PM_v(T)$, the pathosvertices are adjacent to the $v'_1, v'_2, v'_3, \dots, v'_n$ without cut vertex. Thus $PM_v(T)$ is always non-separable.

Case 2. Suppose $T = P_n : v_1, v_2, v_3, \dots, v_n, n > 1$. Further, $V[PM_v(T)] =$

$\{v'_1, v'_2, v'_3 \dots v'_n, e'_1, e'_2, e'_3 \dots e'_{n-1}, r'_1, P'_1\}$. By Theorem 2.1, $M_v(G)$ is non-separable. Also in $PM_v(T)$ the pathosvertex P'_1 adjacent to the vertices $v'_1, v'_2, v'_3 \dots v'_n$ without cut vertex. Hence $PM_v(T)$ is always non-separable.

Theorem 3.2. For any graph G with p vertices, q edges, r regions and k path of pathos, $PM_v(T)$ has $(p + q + r + \sum_{i=1}^k P_i)$ vertices and $q + \sum_{i=1}^q \frac{1}{2}\{d(e_i)\} + \sum_{j=1}^r d(r_j) + \sum_{i=1}^k P_{v_i}$ edges. Where $d(e_i)$ is the edgedegree of a edge e_i , $d(r_j)$ is the degree of a region r_j and P_{v_i} is the number of vertices which lies on the path of pathos.

Proof. By the definition of $PM_v(T)$, $V[PM_v(T)] = V(T)UE(T)UR(T)UP_i(T)$. Hence $V[PM_v(T)] = p + q + r + \sum_{i=1}^k P_i$. Further, by Theorem 2.2, $E[M_v(G)] = q + \sum_{i=1}^q \frac{1}{2}\{d(e_i)\} + \sum_{j=1}^r d(r_j)$. The degree of a pathosvertex is the sum of the number of vertices lies on the each path of pathos in T which is $\sum_{i=1}^k P_{v_i}$. The number of edges in $PM_v(T)$ is equal to the sum of edges in $M_v(G)$ and $\sum_{i=1}^k P_{v_i}$. Hence $E[PM_v(T)] = q + \sum_{i=1}^q \frac{1}{2}\{d(e_i)\} + \sum_{j=1}^r d(r_j) + \sum_{i=1}^k P_{v_i}$.

Theorem 3.3. For any tree T , $PM_v(T)$ is planar if and only if T is P_2 .

Proof. Consider a tree T is not a path P_2 . Let $T = P_3, T = P_3 : v_1, v_2, v_3$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, e'_1, e'_2, r'_1, P'_1\}$. By Theorem 2.3, $M_v(P_3)$ is planar. Further in $PM_v(T)$, the pathos vertex P'_1 is adjacent to the vertices v'_1, v'_2, v'_3 , of $M_v(P_3)$. The edge between v'_2 and P'_1 is crossing over the edge already drawn in $M_v(P_3)$. Which is non-planar, a contradiction. Conversely, suppose $G = P_2 : v_1, v_2$. Further, $V[PM_v(T)] = \{v'_1, v'_2, e'_1, r'_1, P'_1\}$. By Theorem 2.3, $M_v(P_2)$ is planar. Further in $PM_v(T)$, $i[PM_v(P_2)] = 01$. Hence $PM_v(P_2)$ is planar.

Theorem 3.4. For any tree T , $PM_v(T)$ has crossing number one if and only if T is P_3 .

Proof. Consider $T = P_3$, then $M_v(P_3) = 1$ -minimally non outer planar. Further, in $PM_v(P_3)$, the pathos vertex P'_1 is adjacent to the vertices v'_1, v'_2, v'_3 of $M_v(P_3)$. Which gives a crossing number one. Hence $PM_v(T)$ has crossing number one.

Conversely, suppose $T = P_4 : v_1, v_2, v_3, v_4$. Further, $V[M_v(P_4)] = \{v'_1, v'_2, v'_3, v'_4, e'_1, e'_2, e'_3, r'_1, P'_1\}$. By Theorem 2.3, $M_v(P_3)$ is planar. Further in $PM_v(T)$, pathos vertex P'_1 is adjacent to v'_1, v'_2, v'_3, v'_4 and gives crossing number two, a contradiction.

Theorem 3.5. For any tree T , $PM_v(T)$ has crossing number two if and only if T is P_4 or $K_{1,3}$.

Proof. Suppose that $PM_v(T)$ has crossing number two.

We have the following cases.

Case 1. Suppose $T = P_5 : v_1, v_2, v_3, v_4, v_5$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, v'_4, v'_5, e'_1, e'_2, e'_3, e'_4, r'_1, P'_1\}$. By Theorem 2.3, $M_v(P_5)$ is planar. Further in $PM_v(T)$, P'_1 is adjacent to $v'_1, v'_2, v'_3, v'_4, v'_5$ and gives crossing number three, a contradiction.

Case 2. Suppose $T = K_{1,4} : v_1, v_2, v_3, v_4, v_5$ and $deg(v_i) = 4$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, v'_4, v'_5, e'_1, e'_2, e'_3, e'_4, r'_1, P'_1, P'_2\}$. In $M_v(T)$, $C_r[M_v(K_{1,4})] = 06$. Further in $PM_v(T)$, P'_1 adjacent to v'_1, v'_2, v'_3 and P'_2 adjacent to v'_2, v'_4, v'_5 of $M_v(T)$. Which gives a crossing number six, a contradiction.

Conversely,

Case 1. Suppose $T = K_{1,3} : v_1, v_2, v_3, v_4$ and $deg(v_i) = 3$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, v'_4, e'_1, e'_2, e'_3, r'_1, P'_1, P'_2\}$. In $M_v(T)$, $C_r[M_v(K_{1,3})] = 01$. Further in $PM_v(T)$, P'_1 adjacent to v'_1, v'_2, v'_3, v'_4 and P'_2 adjacent to v'_2, v'_3 of $M_v(T)$. Which gives crossing number two. Hence $PM_v(T)$ has crossing number two.

Case 2. Suppose $T = P_4 : v_1, v_2, v_3, v_4$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, v'_4, e'_1, e'_2, e'_3, r'_1, P'_1\}$. By Theorem 2.3, $M_v(P_4)$ is planar. Further in $PM_v(T)$, P'_1 is adjacent to v'_1, v'_2, v'_3, v'_4 and gives crossing number two. Hence $PM_v(T)$ has crossing number two.

Theorem 3.6. For any tree T , $PM_v(T)$ is always noneulerian.

Proof. Let T be any tree. By observation 3.1, the degree of the corresponding vertex in $PM_v(T)$ becomes odd. By the Theorem 2.4, $PM_v(T)$ is noneulerian. Hence, $PM_v(T)$ is always non eulerian.

Theorem 3.7. For any tree T , $PM_v(T)$ is hamiltonian if and only if T is not P_2 .

Proof. Let T be any tree. We have the following cases.

Case 1. Suppose $T = P_n : v_1, v_2, v_3, \dots, v_n, n \geq 3$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, \dots, v'_n, e'_1, e'_2, e'_3, \dots, e'_{n-1}, r'_1, P'_1\}$. Then there exists a cycle $r'_1, v'_1, P'_1, v'_3, e'_2, \dots, v'_n, e'_{n-1}, r'_1$. Which includes all the vertices of $PM_v(T)$. Hence $PM_v(T)$ is hamiltonian.

Case 2. Suppose $T = K_{1,n} : v_1, v_2, v_3, \dots, v_n, n \geq 2$. Further, $V[PM_v(T)] = \{v'_1, v'_2, v'_3, \dots, v'_n, e'_1, e'_2, e'_3, \dots, e'_{n-1}, P'_1, P'_2, \dots, P'_n, r'_1\}$. Then there exists a cycle $r'_1, v'_1, P'_1, v'_3, e'_2, \dots, v'_n, e'_{n-1}, P'_n, r'_1$. Which includes all the vertices of $PM_v(T)$. Hence $PM_v(T)$ is hamiltonian.

Conversely, Suppose $T = P_2 : v_1, v_2$. Further, $V[PM_v(T)] = \{v'_1, v'_2, e'_1, r'_1, P'_1\}$. By the definition of $M_v(T)$, the vertices v'_1 and v'_2 are adjacent to the edge vertex e'_1 and the regionvertex r'_1 . Clearly $M_v(T) = K_4$. In $PM_v(T)$, the pathos vertex P'_1 is adjacent to v'_1 and v'_2 . Clearly $v'_1, e'_1, v'_2, r'_1, v'_1$ form a hamiltonian cycle in $PM_v(T)$. But in $PM_v(T)$, the P'_1 is not lies on the hamiltonian cycle. Hence $PM_v(T)$ is non-hamiltonian.

4. Conclusion. In this paper we obtained the new graph valued function called

pathos vertex semi-middle graph of a tree. We studied the characterization of graphs whose pathos vertex semi-middle graph of a tree is planar, hamiltonian, crossing number one and two. Further, we obtain $PM_v(T)$ is non-eulerian.

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